

Return-into-libc without Function Calls (on the x86)

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Introduction

- ▶ Software-development with C/C++
- ▶ Memory Corruption
 - ▶ Stack overflow
 - ▶ Buffer overflow

Introduction

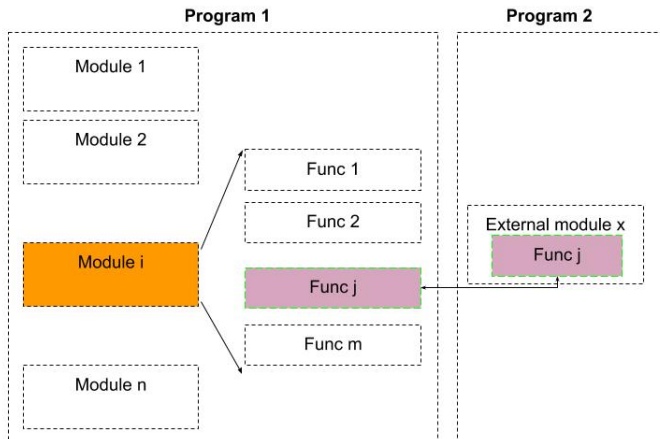
- ▶ Software-development with C/C++
- ▶ Memory Corruption
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- ▶ Code Injection Attacks
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Introduction - Code Injection

- ▶ Function-level
- ▶ External code injecting

Introduction - Code Injection

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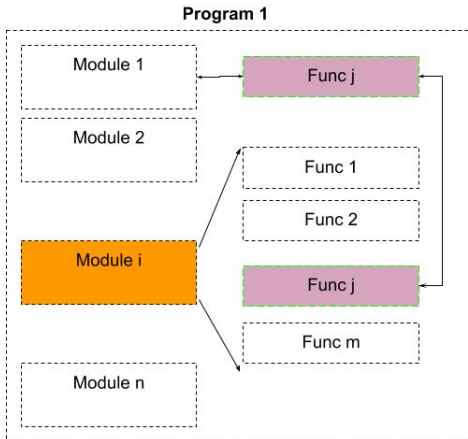


Introduction - Code Reuse

- ▶ Function-level
- ▶ Internal code reuse

Introduction - Code Reuse

- ▶ Function-level
- ▶ Internal code reuse



Problem

- ▶ Removing certain functions from libc
- ▶ Changing the assembler's code generation choices
- ▶ Defense against code reuse attacks

Design

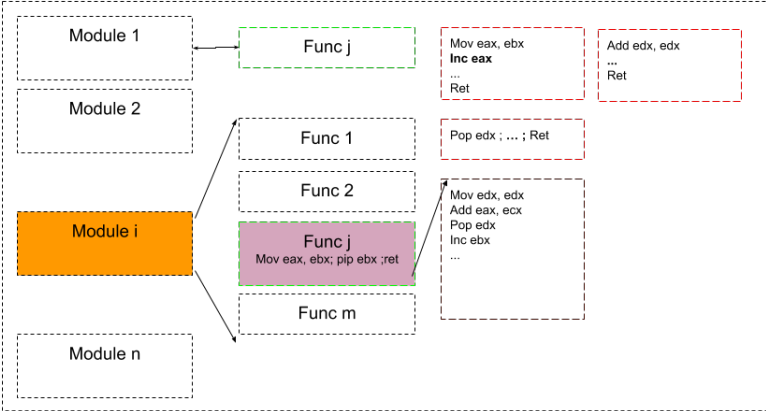
- ▶ Return-oriented programming - ROP
- ▶ Instruction-level - gadgets
- ▶ Discovering useful instructions sequences in Libc

Design

- ▶ Return-oriented programming - ROP
- ▶ Instruction-level - gadgets
- ▶ Discovering useful instructions sequences in Libc
 - ▶ Useful code sequence
 - ▶ Ending with a ret instruction
 - ▶ Boring instructions

Design - ROP

Program 1



Design - GALILEO Algorithm

Algorithm GALILEO:

```
create a node, root, representing the ret instruction;  
place root in the trie;  
for pos from 1 to textseg_len do:  
    if the byte at pos is c3, i.e., a ret instruction, then:  
        call BUILDFROM(pos, root).
```

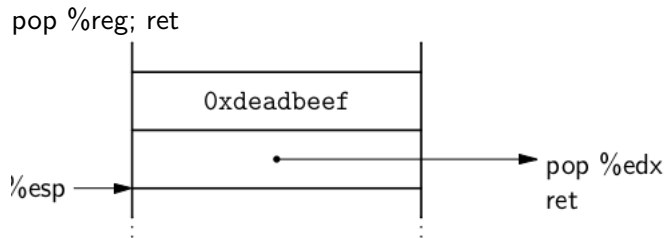
Procedure BUILDFROM(index *pos*, instruction *parent_insn*):

```
for step from 1 to max_insn_len do:  
    if bytes [(pos - step) ... (pos - 1)] decode as a valid instruction insn then:  
        ensure insn is in the trie as a child of parent_insn;  
        if insn isn't boring then:  
            call BUILDFROM(pos - step, insn).
```

Figure 1: The GALILEO Algorithm.

Gadget - Load/Store

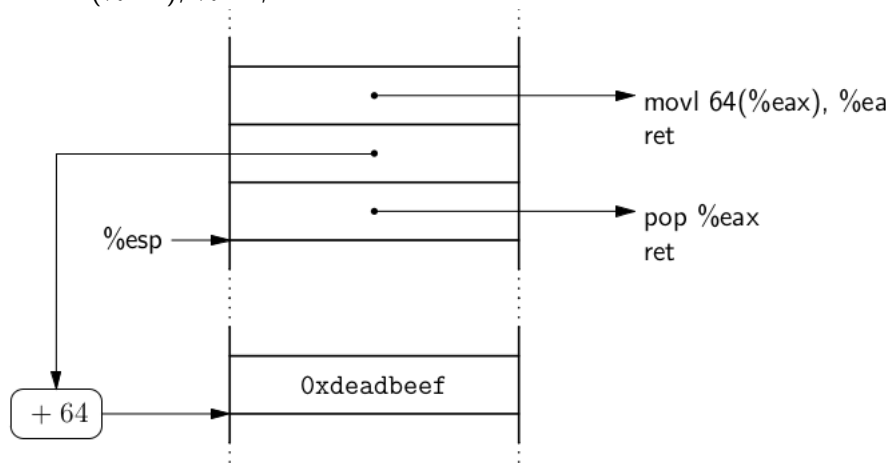
Loading a Constant



Gadget - Load/Store

Loading from Memory

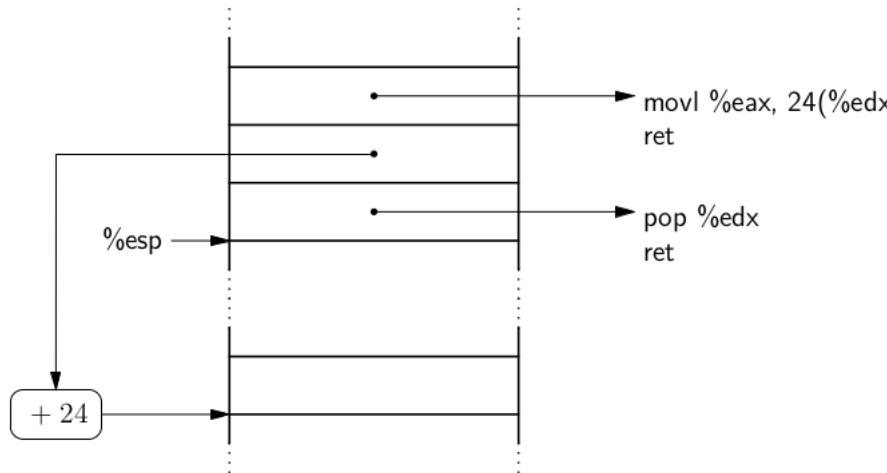
```
movl 64(%eax), %eax; ret
```



Gadget - Load/Store

Storing to Memory

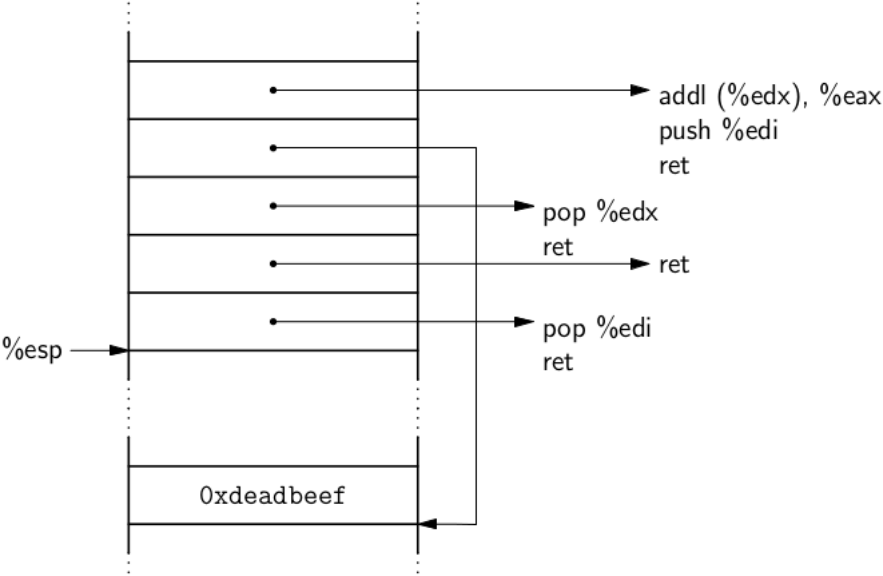
```
movl %eax,24(%edx); ret
```



Gadget - Arithmetic and Logic

Add

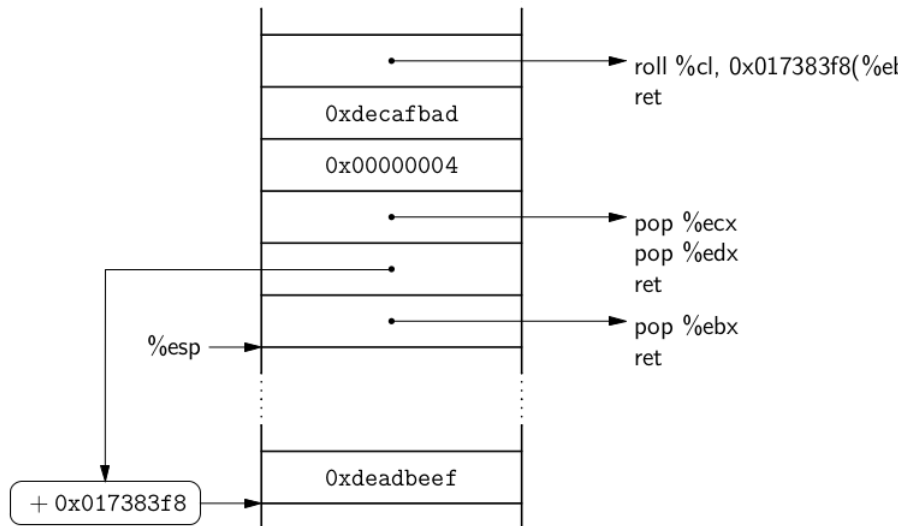
```
addl (%edx),%eax; push %edi; ret
```



Gadget - Arithmetic and Logic

Shifts and Rotates

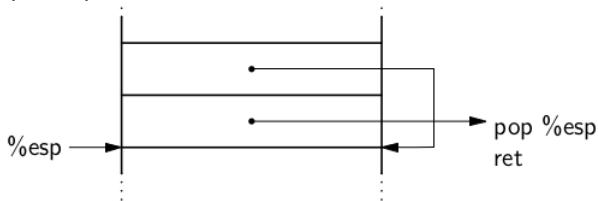
```
roll %cl, 0x17383f8(%ebx);ret
```



Gadget - Control Flow

Unconditional Jumps

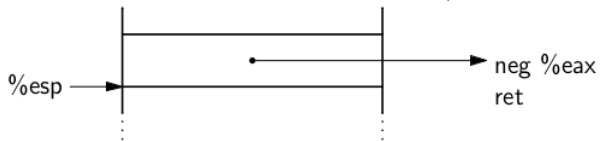
changing the value of `%esp` to point to a new gadget
`pop %esp; ret`



Gadget - Control Flow

Conditional Jumps

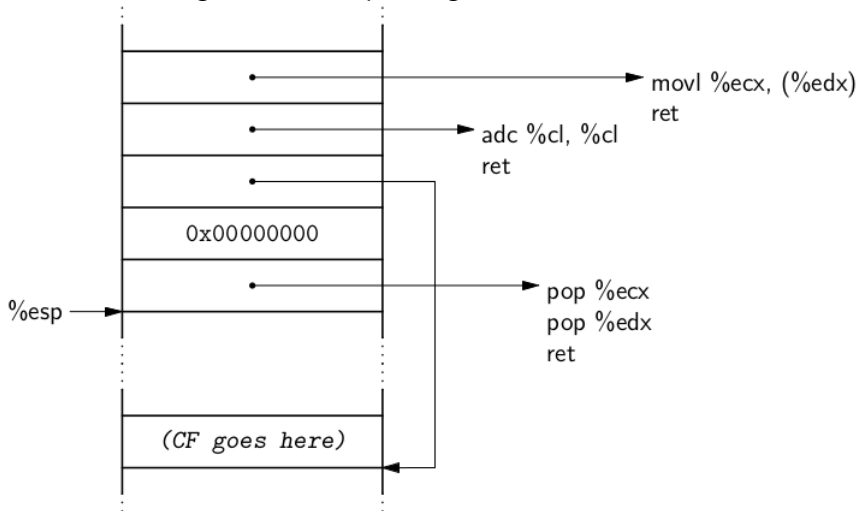
Phase One: Clear CF if %eax is zero, set CF if %eax is nonzero.



Gadget - Control Flow

Conditional Jumps

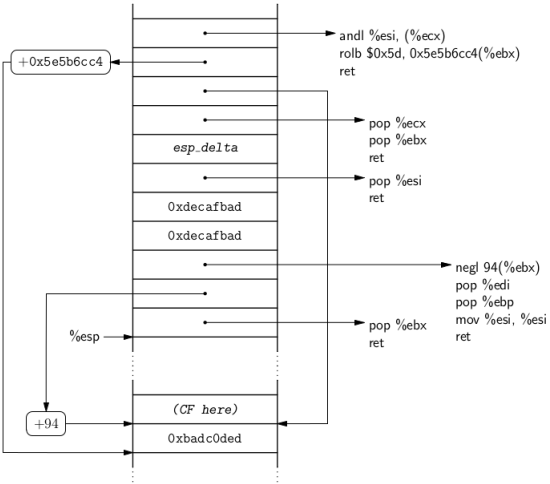
Phase Two: Store either 1 or 0 in the data word labeled “CF goes here,” depending on whether CF is set or not.



Gadget - Control Flow

Conditional Jumps

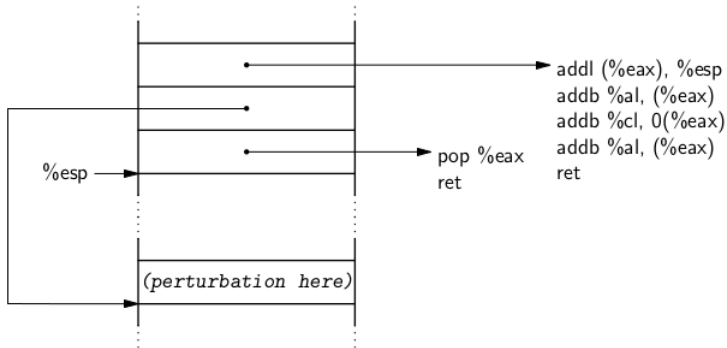
Phase Three: part one: Convert the word (labeled “CF here”) containing either 1 or 0 to contain either esp delta or 0. The data word labeled 0xbadc0ded is used for scratch.



Gadget - Control Flow

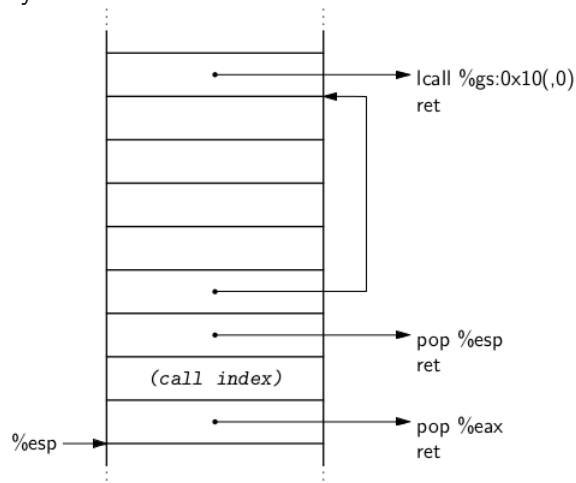
Conditional Jumps

Phase Three: two: Apply the perturbation in the word labeled “perturbation here” to the stack pointer. The perturbation is relative to the end of the gadget.



Gadget - System Calls

System call's number



Implementation

Buffer overflow

- ▶ Buffer overflow vulnerability
- ▶ No randomization
- ▶ No stack-protector

Implementation

Steps

- ▶ `@.data` (@ of `.data` for to place some strings)
- ▶ `int $0x80` (for execute our payload)
- ▶ `mov %eax,(%ecx) — pop %ebp — ret` (for mov eax into buffer)
- ▶ `inc %eax — ret` (for increment eax to up to 11)
- ▶ `pop %edx — pop %ecx — pop %ebx — ret` (for pop address)
- ▶ `pop %eax — pop %ebx — pop %esi — pop %edi — ret` (here just pop `%eax` will be useful)
- ▶ `xor %eax,%eax — ret` (for put `%eax` to zero)

Questions?